

## A Proof of Lemma 3

*Proof of Lemma 3.* We first prove the upper bound of  $A_t$ . The essential proof is actually due to Cheung et al. [2019a] in analyzing sliding window based approach. For self-containedness, we restate here in the notations of our proposed restarted strategy.

$$\begin{aligned} & V_t^{-1} \sum_{s=t_0}^t X_s X_s^T \\ = & V_t^{-1} \sum_{s=t_0}^t X_s X_s^T \sum_{p=s}^{p+1} \\ = & V_t^{-1} \sum_{p=t_0}^t \sum_{s=t_0}^p X_s X_s^T \end{aligned} \quad (21)$$

$$\sum_{p=t_0}^t V_t^{-1} \sum_{s=t_0}^p X_s X_s^T \sum_{p=p+1}^{p+1} \quad (22)$$

$$\sum_{p=t_0}^t \max_{p=p+1} V_t^{-1} \sum_{s=t_0}^p X_s X_s^T k_p \quad (23)$$

$$k_p \quad p+1 k_2; \quad (24)$$

where (21) holds by rearranging over the index pair of  $(s; p)$ , (22) holds due to the triangle inequality, (23) and (24) can be obtained by the same argument in Appendix B of Cheung et al. [2019b]. We thus prove the upper bound of  $A_t$ .

We proceed to prove the upper bound of  $B_t$ . From the self-normalized concentration inequality [Abbasi-Yadkori et al., 2011, Theorem 1], restated in Theorem 5 of Appendix C, we know that

$$\begin{aligned} & \sum_{s=t_0}^t \frac{X_s X_s^T}{V_{t-1}^{-1}} \\ (32) \quad & \leq \frac{2R^2 \log \det(V_{t-1})^{1=2} \det(I)^{-1=2}}{R \left( 2 \log \frac{1}{\epsilon} + d \log \left( 1 + \frac{(t-t_0)L^2}{d} \right) \right)}; \end{aligned}$$

where the last inequality is obtained from the analysis of the determinant, as shown in the proof of Lemma 4.

Meanwhile, since  $V_{t-1} \preceq I$ , we know that

$$k \leq t k_{V_{t-1}^{-1}}^2 = \min(V_{t-1}) k \leq t k_2^2 \leq \frac{1}{k} \leq t k_2^2 \leq S^2;$$

Therefore, the upper bound of  $B_t$  can be immediately obtained by combining the above inequalities.  $\square$

## B Bandit-over-Bandits Mechanism and Proof of Theorem 4

The RestartUCB algorithm requires prior information of the path-length  $P_T$ , which is generally unknown. Such a limitation can be avoided by utilizing the Bandits-over-bandits (BOB) mechanism, proposed by Cheung et al. [2019a] in designing parameter-free algorithm for non-stationary linear bandits based on sliding window least square estimator.

In the following, we first describe how to apply the BOB mechanism to eliminate the requirement of the unknown path-length in RestartUCB. Then, we present the proof of Theorem 4.

### B.1 RestartUCB with BOB Mechanism

We name the RestartUCB algorithm with Bandit-over-Bandits mechanism as "RestartUCB-BOB", whose main idea is illustrated in Figure 4.

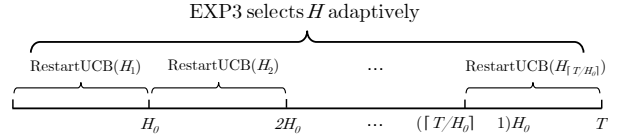


Figure 4: Illustration of Bandit-over-Bandits mechanism with application to RestartUCB algorithm.

From a high-level view, although the exact value of the optimal epoch size (or equivalently, the path-length  $P_T$ ) is not clear, we can make some random guesses of its possible value, since the  $P_T$  is always bounded. Then, we can use a certain meta-algorithm to adaptively track the best epoch size, based on the returned reward returned. Specifically, The RestartUCB-BOB algorithm first sets an update period  $H_0$ , and then runs the RestartUCB with a particular epoch size in each period, and the epoch size will be adaptively adjusted by employing EXP3 [Auer et al., 2002] as the meta-algorithm. We refer the reader to Section 7.3 of Cheung et al. [2019b] for more descriptions of design motivations and algorithmic details.

In the configuration of RestartUCB-BOB, we set  $H_0 = \frac{P_T}{e}$  and the pool of epoch sizes  $J$  as

$$J = \{H_i = b(d=2S)^{2=3} 2^{i-1} c \mid i = 1; 2; \dots; N\};$$

where  $N = \ln(d^{1=3} T^{1=2} (2S)^{2=3}) e + 1$ .

Denoted by  $H_{\min}$  ( $H_{\max}$ ) the minimal (maximal) epoch size in the pool  $J$ , we know that

$$H_{\min} = b(d=2S)^{2=3} c; H_{\max} = b d \frac{P_T}{e} c \quad H_0; \quad (25)$$

## B.2 Proof of Theorem 4

*Proof of Theorem 4.* We begin with the following decomposition of the dynamic regret.

$$\begin{aligned}
 & \sum_{t=1}^T hX_t; t^i - hX_t; t^j \\
 = & \sum_{t=1}^T \underbrace{hX_t; t^i - hX_t(H^y); t^i}_{\text{term (i)}} + \sum_{t=1}^T \underbrace{hX_t(H^y); t^i - X_t(H_i); t^i}_{\text{term (ii)}}
 \end{aligned}$$

where  $H^y$  is the best epoch size to approximate the optimal epoch size  $H$  in the pool  $J$ , and  $H = b(dT/(1+P_T))^{2/3}c$ . Hence, it suffices to bound terms (i) and (ii). In the following, we consider two cases, either  $(1+P_T) \geq d^{1/2}T^{1/4}$  or  $(1+P_T) < d^{1/2}T^{1/4}$ .

**Case 1.** when  $(1+P_T) \geq d^{1/2}T^{1/4}$ .

In this case, it is easy to verify that  $H \leq H_{\max}$  and we thus conclude that  $H$  lies in the range of  $[H_{\min}; H_{\max}]$ . Furthermore, from the configuration of the pool  $J$ , we confirm that there exists an epoch size  $H^y \geq J$  such that  $H^y \leq H \leq 2H^y$ . So term (ii) can be upper bounded by

$$\text{term (ii)} \leq \sum_{i=1}^{dT/H_0} H^y P_i + \frac{dH_0}{H^y} \quad (26)$$

$$\begin{aligned}
 & = \Theta(H^y P_T + \frac{dT}{H^y}) \quad (27) \\
 & \leq \Theta(H P_T + \frac{dT}{2H}) \\
 & = \Theta(d^{2/3} P_T^{1/3} T^{2/3});
 \end{aligned}$$

where (26) is due to Theorem 2 and  $P_i$  denotes the path-length in the  $i$ -th update period. (27) follows by summing over all update periods, and the last inequality holds since the optimal epoch size  $H$  is provably in the range of  $[H_{\min}; H_{\max}]$  and satisfies  $H^y \leq H \leq 2H^y$ .

Next, we bound the term (i),

$$\begin{aligned}
 \text{term (i)} & \leq \frac{dT}{H_0 N T} \\
 & \leq \Theta(d^{1/2} T^{3/4}) \quad (28) \\
 & \leq \Theta(d^{2/3} T^{2/3} (1+P_T)^{1/3});
 \end{aligned}$$

where the first inequality follows by the same argument as in the sliding window based approach [Cheung et al., 2019b, Lemma 13], building upon the of EXP3. In

addition, the last inequality holds due to the fact that  $(1+P_T) \geq d^{1/2}T^{1/4}$  implies,

$$d^{1/2} T^{3/4} = d^{2/3} T^{2/3} d^{1/3} T^{1/6} \leq d^{2/3} T^{2/3} (1+P_T)^{1/3};$$

Hence, by combining the upper bounds of term (i) and term (ii), we know that the dynamic regret of RestartUCB-BOB is bounded by  $\Theta(d^{2/3} T^{2/3} (1+P_T)^{1/3})$  under the condition of  $(1+P_T) \geq d^{1/2}T^{1/4}$ .

**Case 2.** when  $(1+P_T) < d^{1/2}T^{1/4}$ .

In this case, we cannot guarantee that the optimal epoch size  $H$  lies in the range of  $[H_{\min}; H_{\max}]$ , so we set  $H^y$  as  $H_0$ ,

$$\begin{aligned}
 \text{term (i)} & \leq H^y P_T + \frac{dT}{H^y} \\
 & \leq H_0 P_T + \frac{dT}{H_0} \\
 & = \Theta(d^{1/2} T^{1/4} P_T + d^{1/2} T^{3/4}) \\
 & \leq \Theta(d^{1/2} T^{3/4})
 \end{aligned}$$

where the last inequality holds by exploiting the condition of  $(1+P_T) \geq d^{1/2}T^{1/4}$ . The result in conjunction with the upper bound of term (i) in (28) gives the  $\Theta(d^{1/2} T^{3/4})$  dynamic regret under this condition.

Finally, note that the dynamic regret of above two cases can be rewritten in the following unified form,

$$\text{term (i)} + \text{term (ii)} \leq \Theta(d^{2/3} T^{2/3} \max\{P_T; d^{1/2} T^{1/4}\} g^{1/3});$$

Hence, we complete the proof of Theorem 4.  $\square$

## C Technical Lemmas

In this section, we provide several technical lemmas that frequently used in the proofs.

**Theorem 5** (Self-Normalized Bound for Vector-Valued Martingales [Abbasi-Yadkori et al., 2011, Theorem 1]). *Let  $f: F_t \rightarrow \mathbb{R}^d$  be a filtration. Let  $f: g_{t=0}^1$  be a real-valued stochastic process such that  $f_t$  is  $F_t$ -measurable and conditionally  $R$ -sub-Gaussian for some  $R > 0$ , namely,*

$$\mathbb{E}[\exp(-f_t) | F_{t-1}] \leq \exp\left(-\frac{f_t^2 R^2}{2}\right); \quad (29)$$

*Let  $f: X_t: g_{t=1}^1$  be an  $\mathbb{R}^d$ -valued stochastic process such that  $X_t$  is  $F_{t-1}$ -measurable. Assume that  $V$  is a  $d \times d$  positive definite matrix. For any  $t \geq 0$ , define*

$$V_t = V + \sum_{s=1}^t X_s X_s^T; \quad S_t = \sum_{s=1}^t X_s; \quad (30)$$

Then, for any  $\epsilon > 0$ , with probability at least  $1 - \epsilon$ , for all  $t \geq 0$ ,

$$kS_t k_{V_t}^2 \leq 2R^2 \log \frac{\det(V_t)^{1-2} \det(V_0)^{1-2}}{\det(V_0)}; \quad (31)$$

**Lemma 4** (Elliptical Potential Lemma). Suppose  $U_0 = I$ ,  $U_t = U_{t-1} + X_t X_t^T$ , and  $kX_t k_2 \leq L$ , then

$$\sum_{t=1}^T kU_t^{-\frac{1}{2}} X_t k_2 \leq \sqrt{2dT \log \left(1 + \frac{L^2 T}{d}\right)}; \quad (32)$$

*Proof.* First, we have the following decomposition,

$$U_t = U_{t-1} + X_t X_t^T = U_{t-1}^{\frac{1}{2}} \left( I + U_{t-1}^{-\frac{1}{2}} X_t X_t^T U_{t-1}^{-\frac{1}{2}} \right) U_{t-1}^{\frac{1}{2}};$$

Taking the determinant on both sides, we get

$$\det(U_t) = \det(U_{t-1}) \det\left(I + U_{t-1}^{-\frac{1}{2}} X_t X_t^T U_{t-1}^{-\frac{1}{2}}\right);$$

which in conjunction with Lemma 5 yields

$$\begin{aligned} \det(U_t) &= \det(U_{t-1}) \left(1 + kU_{t-1}^{-\frac{1}{2}} X_t k_2^2\right) \\ &= \det(U_{t-1}) \exp(kU_{t-1}^{-\frac{1}{2}} X_t k_2^2); \end{aligned}$$

Note that in the first inequality, we utilize the fact that  $1 + x \leq \exp(x)$  holds for any  $x \geq 0$ . By taking advantage of the telescope structure, we have

$$\sum_{t=1}^T kU_t^{-\frac{1}{2}} X_t k_2 \leq 2 \log \frac{\det(U_T)}{\det(U_0)} \leq 2d \log \left(1 + \frac{L^2 T}{d}\right);$$

where the last inequality follows from the fact that  $\text{Tr}(U_T) = \text{Tr}(U_0) + L^2 T = d + L^2 T$ , and thus  $\det(U_T) \leq (d + L^2 T)^d$ .

Therefore, Cauchy-Schwartz inequality gives,

$$\sum_{t=1}^T kU_t^{-\frac{1}{2}} X_t k_2 \leq \sqrt{\sum_{t=1}^T kU_t^{-\frac{1}{2}} X_t k_2^2} \leq \sqrt{2dT \log \left(1 + \frac{L^2 T}{d}\right)};$$

□

**Lemma 5.**

$$\det(I + \mathbf{v}\mathbf{v}^T) = 1 + k\mathbf{v} k_2^2; \quad (33)$$

*Proof.* Notice that

- (i)  $(I + \mathbf{v}\mathbf{v}^T)\mathbf{v} = (1 + k\mathbf{v} k_2^2)\mathbf{v}$ , therefore,  $\mathbf{v}$  is its eigenvector with  $(1 + k\mathbf{v} k_2^2)$  as the eigenvalue;
- (ii)  $(I + \mathbf{v}\mathbf{v}^T)\mathbf{v}^\perp = \mathbf{v}^\perp$ , therefore,  $\mathbf{v}^\perp$  is its eigenvector with 1 as the eigenvalue.

Consequently,  $\det(I + \mathbf{v}\mathbf{v}^T) = 1 + k\mathbf{v} k_2^2$ . □